A few constructive approaches to optimal first-order optimization methods for convex optimization

Adrien Taylor





All Russian optimization seminar - May 2021

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If interested, details are provided in references at the end. Complementary material on Francis Bach's blog (also \pm informal) https://francisbach.com/computer-aided-analyses/

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But also:

- ◇ Fair amount of algorithmic analyses (and design) originated from SDPs (from different authors, examples below), in different settings.
- ◊ We try keeping track of related works in the toolbox' manual (see later), incomplete references in this presentation.



François Glineur



Yoel Drori



Julien Hendrickx



Francis Bach



Etienne de Klerk



Jérôme Bolte



Van Scoy



Ernest Ryu



Alexandre d'Aspremont



Mathieu Barré



Radu-Alexandru Dragomir

Carolina

Bergeling





Pontus Giselsson



Laurent Lessard

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When it does: principled approach to worst-case analyses.

Performance estimation problems

Designing methods using PEPs

Conclusions

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Question: what a priori guarantees after N iterations?

Examples: how small should $f(x_N) - f(x_*)$, $||f'(x_N)||$, $||x_N - x_*||$ be?

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In other words:

$$\begin{split} \|f'(x_N)\| &\leqslant \sup_{F, y_0, \dots, y_N} \quad \|F'(y_N)\| \\ & \text{subject to} \quad y_1, \dots, y_N \text{ generated by gradient method from } y_0 \\ & F \text{ satisfies the assumptions on } f \end{split}$$

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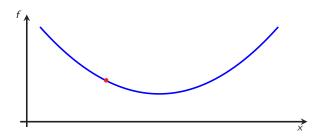
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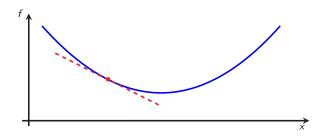
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Smooth strongly convex functions

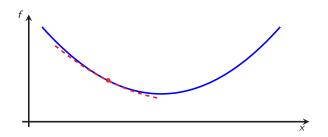
Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:



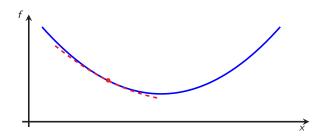
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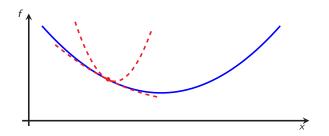


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Toy example: What can we guarantee on $||f'(x_1)||$ given that:

- $\diamond \ f$ is L-smooth and μ -strongly convex (notation $f \in \mathcal{F}_{\mu,\mathsf{L}}$),
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 Require points (x_i, g_i, f_i) to be interpolable by a function f ∈ F_{µ,L}. The new constraint is:

$$\exists f \in \mathcal{F}_{\mu,L}: f_i = f(x_i), g_i = f'(x_i), \quad \forall i \in \{0,1\}.$$

◊ Performance estimation problem:

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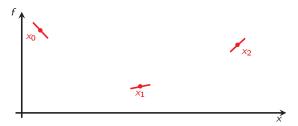
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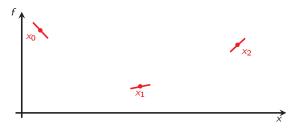
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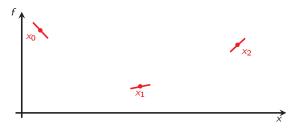
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$$f_i \geq f_j + \langle g_j, x_i - x_j \rangle + \frac{1}{2L} \|g_i - g_j\|^2 + \frac{\mu}{2(1-\mu/L)} \|x_i - x_j - \frac{1}{L}(g_i - g_j)\|^2.$$

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- Simpler example: pick $\mu = 0$ and $L = \infty$ (just convexity):

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$$\begin{split} f_1 &\ge f_0 + \langle g_0, x_1 - x_0 \rangle + \frac{1}{2L} \|g_1 - g_0\|^2 + \frac{\mu}{2(1 - \mu/L)} \|x_1 - x_0 - \frac{1}{L} (g_1 - g_0)\|^2 \\ f_0 &\ge f_1 + \langle g_1, x_0 - x_1 \rangle + \frac{1}{2L} \|g_0 - g_1\|^2 + \frac{\mu}{2(1 - \mu/L)} \|x_0 - x_1 - \frac{1}{L} (g_0 - g_1)\|^2. \end{split}$$

♦ Same optimal value (no relaxation); but still non-convex quadratic problem.

♦ Using $x_1 = x_0 - \gamma g_0$, all elements are quadratic in (g_0, g_1) , and linear in (f_0, f_1) :

$$\begin{split} \max_{\substack{g_0,g_1\\f_0,f_1}} & \|g_1\|^2 \\ \text{subject to} & f_1 \geqslant f_0 - \gamma \|g_0\|^2 + \frac{1}{2L} \|g_1 - g_0\|^2 + \frac{\mu}{2(1-\mu/L)} \left\| \gamma g_0 + \frac{1}{L} (g_1 - g_0) \right\|^2 \\ & f_0 \geqslant f_1 + \gamma \langle g_1, g_0 \rangle + \frac{1}{2L} \|g_1 - g_0\|^2 + \frac{\mu}{2(1-\mu/L)} \left\| \gamma g_0 + \frac{1}{L} (g_1 - g_0) \right\|^2 \\ & \|g_0\|^2 = R^2. \end{split}$$

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 \diamond They are therefore linear in terms of a Gram matrix G and a vector F, with

$$G = \begin{bmatrix} \|g_0\|^2 & \langle g_0, g_1 \rangle \\ \langle g_0, g_1 \rangle & \|g_1\|^2 \end{bmatrix} = \begin{bmatrix} g_0 & g_1 \end{bmatrix}^\top \begin{bmatrix} g_0 & g_1 \end{bmatrix}, \quad F = \begin{bmatrix} f_0 & f_1 \end{bmatrix},$$

where $G \geq 0$ by construction.

 \diamond Using the new variables $G \succeq 0$ and F

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 $\diamond~$ previous problem can be reformulated as a 2 \times 2 SDP

$$\begin{array}{ll} \max_{G,F} & G_{2,2} \\ \text{subject to} & F_1 - F_0 + \frac{\gamma L(2 - \gamma \mu) - 1}{2(L - \mu)} G_{1,1} + \frac{1 - \gamma \mu}{L - \mu} G_{1,2} - \frac{1}{2(L - \mu)} G_{2,2} \ge 0 \\ & F_0 - F_1 + \frac{\gamma \mu (2 - \gamma L) - 1}{2(L - \mu)} G_{1,1} + \frac{1 - \gamma L}{L - \mu} G_{1,2} - \frac{1}{2(L - \mu)} G_{2,2} \ge 0 \\ & G_{1,1} = 1 \\ & G \succcurlyeq 0. \end{array}$$

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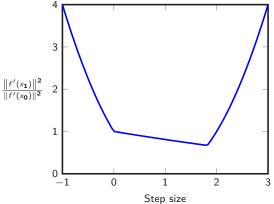
♦ For d = 1 same optimal value by adding rank(G) ≤ 1 .

Solving the SDP...

Fix L = 1, $\mu = .1$ and solve the SDP for a few values of γ .

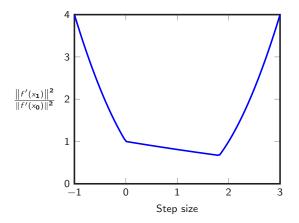
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Observation: numerics match the (expected) $\max\{(1 - \gamma L)^2, (1 - \gamma \mu)^2\}$.

 \diamond Let us rephrase our target: we look for $\rho(\gamma)$ (hopefully small) such that

 $\left\|f'(x_1)\right\| \le \rho(\gamma) \left\|f'(x_0)\right\|$

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- \diamond Feasible points to the previous SDP correspond to lower bounds on $\rho(\gamma)$.
- ♦ Any such $\rho(\gamma)$ that is valid for all $x_0 \in \mathbb{R}^d$, $d \in \mathbb{N}$, $f \in \mathcal{F}_{\mu, \mathsf{L}}$ is a feasible point to the dual SDP.

◊ Recall primal problem

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- $\diamond~$ Introduce dual variables $\lambda_1,~\lambda_2$ and τ for the linear constraints, and dualize.
- ◊ Dual problem is

$$\begin{array}{l} \underset{\tau,\lambda_{1},\lambda_{2} \geqslant 0}{\text{minimize } \tau} \\ \text{subject to } S = \begin{bmatrix} -\frac{\lambda_{1}(\gamma\mu-1)(\gamma L-1)}{L-\mu} - \tau & -\frac{\lambda_{1}(\gamma(\mu+L)-2)}{2(L-\mu)} \\ -\frac{\lambda_{1}(\gamma(\mu+L)-2)}{2(L-\mu)} & 1 - \frac{\lambda_{1}}{L-\mu} \end{bmatrix} \preccurlyeq 0 \\ 0 = \lambda_{1} - \lambda_{2}. \end{array}$$

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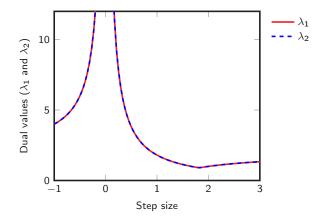
◇ Strong duality holds (existence of a Slater point): any valid worst-case convergence rate ≡ valid dual feasible point (↓) : hence "↑".

Solving the dual

Fix L = 1, $\mu = .1$ and solve the dual SDP for a few values of γ .

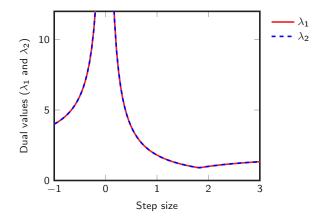
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Note: numerics match $\lambda_1 = \lambda_2 = \frac{2}{|\gamma|}\rho(\gamma)$ with $\rho(\gamma) = \max\{|1 - \gamma L|, |1 - \gamma \mu|\}.$

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$$\begin{split} f_{0} &\geq f_{1} \quad + \langle f'(x_{1}), x_{0} - x_{1} \rangle + \frac{1}{2L} \| f'(x_{0}) - f'(x_{1}) \|^{2} \\ &+ \frac{\mu}{2(1 - \mu/L)} \| x_{0} - x_{1} - \frac{1}{L} (f'(x_{0}) - f'(x_{1})) \|^{2} \\ f_{1} &\geq f_{0} \quad + \langle f'(x_{0}), x_{1} - x_{0} \rangle + \frac{1}{2L} \| f'(x_{0}) - f'(x_{1}) \|^{2} \\ &+ \frac{\mu}{2(1 - \mu/L)} \| x_{0} - x_{1} - \frac{1}{L} (f'(x_{0}) - f'(x_{1})) \|^{2} \\ \end{split} : \lambda_{2} = \frac{2}{\gamma} (1 - \mu\gamma)$$

with $\lambda_1, \lambda_2 \ge 0$. Weighted sum can be reformulated as

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leading to $\|f'(x_1)\|^2 \leqslant (1-\frac{\mu}{L})^2 \|f'(x_0)\|^2$ (tight).

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For finding proofs:

- ◊ the SDP might help by playing with both sides:
 - play with primal (e.g., worst-case functions might be easy to identify),
 - play with dual (e.g., dual variables might be easy to identify).
- ◊ Standard tricks apply, e.g., trace norm minimization for promoting low-rank solutions (on primal or dual).

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- \diamond pick a method
- ◊ pick a class of functions
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- ... such conditions (or slight generalizations) hold in variety of cases (see later).

In other situations, one might want to relax the PEP for obtaining upper-bounds.

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- ◊ Optimizing/designing methods? upcoming!

Avoiding semidefinite programming modeling steps?

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François Glineur (UCLouvain)



Julien Hendrickx (UCLouvain)

"Performance Estimation Toolbox (PESTO): automated worst-case analysis of first-order optimization methods" (CDC 2017)

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Fast Gradient Method (FGM) Input: $f \in \mathcal{F}_{0,L}(\mathbb{R}^d)$, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L}f'(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2}(x_{i+1} - x_i)$

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What if inexact gradient used instead? Relative inaccuracy model:

$$\|\widetilde{d}_f(y_i) - f'(y_i)\| \leq \varepsilon \|f'(y_i)\|.$$

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```
% (0) Initialize an empty PEP
P = pep();
```

% (1) Set up the	objective function
param.mu = 0;	% strong convexity parameter
param.L = 1;	% Smoothness parameter

F=P.DeclareFunction('SmoothStronglyConvex', param); % F is the objective function

```
% (2) Set up the starting point and initial condition

x0 = P.StartingPoint(); % x0 is some starting point

[xs, fs] = F.OptimalPoint(); % xs is an optimal point, and fs=F(xs)

P.InitialCondition((x0-xs)^2 <= 1); % Add an initial condition ||x0-xs||^2<= 1
```

```
% (3) Algorithm
N = 7: % number of iterations
```

```
x = cell(N=1,1); % we store the iterates in a cell for convenience
x(1) = x0;
y = x0;
eps = .1;
for i = 1:N
d = inexactsubgradient(y, F, eps);
x(i+1) = y - 1/param.L * d;
y = x(i+1) + (i-1)/(i+2) * (x(i+1) - x(i));
end
```

```
% (4) Set up the performance measure
[g, f] = F.oracle(x{N+1}); % g=grad F(x), f=F(x)
P.PerformanceMetric(f - fs); % Worst-case evaluated as F(x)-F(xs)
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% (5) Solve the PEP 
P.solve()
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% (6) Evaluate the output
double(f - fs) % worst-case objective function accuracy
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0.3

= 0.1= 0.0

PESTO example: Douglas-Rachford splitting

```
% (0) Initialize an empty PEP
 P=pep():
 N = 1:
% (1) Set up the class of monotone inclusions
paramA.L = 1; paramA.mu = 0; % A is 1-Lipschitz and 0-strongly monotone
paramB.mu = .1; % B is .1-strongly monotone
 A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
 B = P.DeclareFunction('StronglyMonotone', paramB);
w = cell(N+1,1); wp = cell(N+1.1):
x = cell(N, 1); xp = cell(N, 1);
y = cell(N,1); yp = cell(N.1);
% (2) Set up the starting points
w{1} = P.StartingPoint(): wp{1} = P.StartingPoint():
 P.InitialCondition((will-wpill)^2<=1):
% (3) Algorithm
lambda = 1.3: % step size (in the resolvents)
 theta = .9: % overrelaxation
If k = 1 : N
     x{k} = proximal step(w{k}.B.lambda):
     y{k} = proximal step(2*x{k}-w{k},A,lambda);
     w\{k+1\} = w\{k\} \cdot theta*(x\{k\} \cdot v\{k\}):
     xp{k} = proximal step(wp{k}.B.lambda);
     yp{k} = proximal step(2*xp{k}-wp{k},A,lambda);
     wp\{k+1\} = wp\{k\} \cdot theta*(xp\{k\} \cdot vp\{k\});
- end
% (4) Set up the performance measure: ||z0-z1||^2
 P.PerformanceMetric((w{k+1}-wp{k+1})^2):
 % (5) Solve the PEP
 P.solve()
 % (6) Evaluate the output
 double((w{k+1}-wp{k+1})^2) % worst-case contraction factor
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B = P.DeclareFunction('StronglyMonotone', paramB);
w = cell(N+1,1); wp = cell(N+1.1):
x = cell(N, 1); xp = cell(N, 1);
v = cell(N, 1); v_D = cell(N, 1);
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 x{k]
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                                                                ρ3
                                                                    0.8
A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
                                                                 Contraction rate
B = P.DeclareFunction('StronglyMonotone', paramB);
                                                                    0.6
w = cell(N+1,1);
                   wp = cell(N+1,1);
x = cell(N, 1);
                    xp = cell(N, 1);
                                                                    0.4
v = cell(N, 1):
                    vp = cell(N, 1):
                                                                    0.2
% (2) Set up the starting points
w{1} = P.StartingPoint(): wp{1} = P.StartingPoint():
P.InitialCondition((will-wpill)^2<=1):
                                                                              0.5
                                                                                      1
                                                                                            1.5
                                                                            Lipschitz constant L
% (3) Algorithm
lambda = 1.3:
                    % step size (in the resolvents)
theta = .9:
                    % overrelaxation
 x{k}
            = proximal step(w{k},B,lambda);
            = proximal step(2*x{k}-w{k},A,lambda);
 v{k}
 w{k+1}
            = w{k}-theta*(x{k}-y{k});
            = proximal step(wp{k}.B.lambda):
    xp{k}
    vp{k}
            = proximal step(2*xp{k}-wp{k},A,lambda);
            = wp{k}-theta*(xp{k}-vp{k}):
    wp{k+1}
- end
% (4) Set up the performance measure: ||z0-z1||^2
P.PerformanceMetric((w{k+1}-wp{k+1})^2):
% (5) Solve the PEP
P.solve()
% (6) Evaluate the output
```

 $\mu = 0.1$

 $\mu = 0.5$

 $\mu = 1.5$

 $\mu = 1$

 $\mu = 2$

2

PESTO example: Douglas-Rachford splitting

```
% (0) Initialize an empty PEP
P=pep():
N = 1:
% (1) Set up the class of monotone inclusions
paramA.L = 1: paramA.mu = 0; % A is 1-Lipschitz and 0-strongly monotone
paramB.mu = .1;
                              % B is .1-strongly monotone
                                                                °
                                                                                                              \mu = 0.1
                                                                   0.8
A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
                                                                Contraction rate
                                                                                                              \mu = 0.5
B = P.DeclareFunction('StronglyMonotone', paramB);
                                                                   0.6
                                                                                                              \mu = 1
                   wp = cell(N+1,1);
w = cell(N+1.1):
                                                                                                              \mu = 1.5
x = cell(N, 1);
                    xp = cell(N, 1);
                                                                   0.4
v = cell(N, 1):
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                                                                                                              \mu = 2
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                                                                             0.5
                                                                                     1
                                                                                           1.5
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    xp{k}
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    vp{k}
            = proximal step(2*xp{k}-wp{k},A,lambda);
             = wp{k} - theta*(xp{k} - yp{k});
    wp{k+1}
- end
                                                      \checkmark fast prototyping (~ 20 effective lines)
% (4) Set up the performance measure: ||z0-z1||^2
                                                      \checkmark quick analyses (\sim 10 minutes)
P.PerformanceMetric((w{k+1}-wp{k+1})^2):

    computer-aided proofs (multipliers)

% (5) Solve the PEP
P.solve()
% (6) Evaluate the output
double((w{k+1}-wp{k+1})^2)
                            % worst-case contraction factor
                                                                                                                29
```

Includes... but not limited to

- ◊ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- proximal point algorithm,
- o projected and proximal gradient, accelerated/momentum versions,
- ◊ steepest descent, greedy/conjugate gradient methods,
- Douglas-Rachford/three operator splitting,
- ◊ Frank-Wolfe/conditional gradient,
- ◊ inexact gradient/fast gradient,
- ◊ Krasnoselskii-Mann and Halpern fixed-point iterations,
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Currently in Matlab, soon in Python.

Performance estimation problems

Designing methods using PEPs

Conclusions

Main inspiration

Great inspiration from previous works.

- ♦ B. Polyak. "Introduction to optimization" (1964)
- $\diamond~$ Y. Nesterov. "A method of solving a convex programming problem with convergence rate $O(1/k^2)$ ". (1983)
- ◊ A. Nemirovsky, and B. Polyak. "Iterative methods for solving linear ill-posed problems under precise information." (1984)
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In next couple of slides:

- ◊ goal: principled way towards optimal methods.
- ◊ in some sense, generalization of Chebyshev methods (tailored for quadratic minimization) to non-quadratic smooth strongly convex setup.



Yoel Drori



Main references for the following slides (sloppy references throughout):

- ◊ Y. Drori, T., "On the oracle complexity of smooth strongly convex minimization". (2021)
- T., Y. Drori, "An optimal gradient method for smooth strongly convex minimization". (2021)
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- ◊ Y. Drori, "The exact information-based complexity of smooth convex minimization". (2017)
- ♦ D. Kim, J.F. Fessler, "Optimized first-order methods for smooth convex minimization". (2016)
- ◊ Y. Drori, M. Teboulle, "Performance of first-order methods for smooth convex minimization: a novel approach". (2014)



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- ◊ Y. Drori, M. Teboulle, "Performance of first-order methods for smooth convex minimization: a novel approach". (2014)

Also closely related:

- ♦ B. Van Scoy, R.A. Freeman, K.M. Lynch, "The fastest known globally convergent first-order method for minimizing strongly convex functions". (2017)
- ◊ D. Kim, J.F. Fessler, "Optimizing the efficiency of first-order methods for decreasing the gradient of smooth convex functions". (2021)

How could we use PEPs for designing methods, and lower complexity bounds?

 $\diamond~$ Upper bound side: PEP is a machinery for designing worst-case guarantees

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 - provide a constructive way to generate worst-case examples,
 - can be used for designing "worst functions in the world".

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(iii) A notion of "accuracy". Examples: $\frac{f(w_N)-f_\star}{\|w_0-w_\star\|^2}$, $\frac{\|w_N-w_\star\|^2}{\|w_0-w_\star\|^2}$, $\frac{\|f'(w_N)\|^2}{f(w_0)-f_\star}$. \Rightarrow Resulting design problem, for example

$$\min_{\{h_{i,j}\}} \max_{f \in \mathcal{F}_{\mathbf{0},L}} \left\{ \frac{f(w_N) - f_{\star}}{\|w_0 - w_{\star}\|^2} : w_N \text{ obtained from (FOM) and } w_0 \right\}.$$

(i.e., "minimize worst-case")

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– Ex: for
$$\frac{\|w_N - w_\star\|^2}{\|w_0 - w_\star\|^2}$$
, we get "Chebyshev semi-iterative method":

$$w_k = w_{k-1} - \frac{4\delta_k}{L-\mu} f'(w_{k-1}) + \left(1 - \frac{2\delta_k(L+\mu)}{L-\mu}\right) (w_{k-2} - w_{k-1}),$$

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♦ The situation actually quite similar beyond quadratics.

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 - those methods are incredibly close to Nesterov's method!
- Beyond that, a few criterion/settings/methods for which "perfectly optimal" algorithms might be known, but matching lower bounds are still missing.

$$-\frac{\|f'(w_N)\|^2}{f(w_0)-f_1}$$
,

a few numerically-generated methods.

 $\begin{array}{l} \label{eq:optimized gradient method} (OGM) \\ \text{Example I: the optimal method for } \frac{f(y_N) - f_*}{\|y_0 - y_*\|^2} \text{ is called the "optimized gradient} \\ \text{method" (OGM, Kim \& Fessler (2016)):} \end{array}$

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$$y_k = \frac{1}{\theta_{k,N}} z_k + \left(1 - \frac{1}{\theta_{k,N}}\right) \left(y_{k-1} - \frac{1}{L} f'(y_{k-1})\right)$$
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which rely on some exotic sequence

$$\theta_{k+1,N} = \begin{cases} \frac{1 + \sqrt{4\theta_{k,N}^2 + 1}}{2} & \text{if } k \leqslant N - 2\\ \frac{1 + \sqrt{8\theta_{k,N}^2 + 1}}{2} & \text{if } k = N - 1, \end{cases}$$

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- ◊ Nesterov's method: can be obtained as an optimized gradient method whose proof relies only on more convenient inequalities.

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with step sizes

$$[h_{i,j}^{\star}] = \begin{bmatrix} 1.9060 \\ 0.3879 & 2.1439 \\ 0.1585 & 0.4673 & 2.1227 \\ 0.0660 & 0.1945 & 0.4673 & 2.1439 \\ 0.0224 & 0.0660 & 0.1585 & 0.3879 & 1.9060 \end{bmatrix}$$

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$$[h_{i,j}^{\star}] = \begin{bmatrix} 1.9470 \\ 0.4599 & 2.2406 \\ 0.1705 & 0.4599 & 1.9470 \end{bmatrix}$$

Shape of lower complexity bounds

Role of extension/interpolation results, so far?

- ♦ For obtaining *tight* SDP representation of the worst-case computation problem.
- $\diamond~$ We can infer shapes for the worst-case functions!
 - Why? Let's flashback into the interpolation/extension problem!

Reminder: smooth strongly convex interpolation/extension

Consider a set S, and its associated values $\{(x_i, g_i, f_i)\}_{i \in S}$ with coordinates x_i , subgradients g_i and function values f_i .

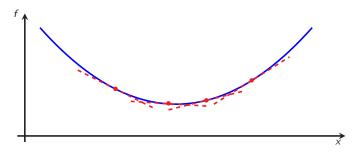


? Possible to find a $f \in \mathcal{F}_{\mu,L}$ s.t.

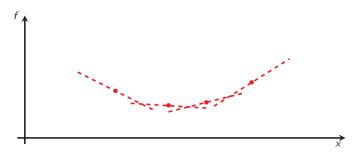
 $f(x_i) = f_i$, and $g_i \in \partial f(x_i)$, $\forall i \in S$.

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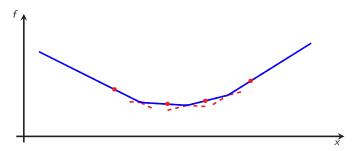


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Explicit construction:

$$f(x) = \max_{j} \left\{ f_{j} + \left\langle g_{j}, x - x_{j} \right\rangle \right\},\,$$

Not unique.

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 - such functions are sometimes referred to as being "zero-chain".

Worst-case performance $\frac{f(w_k) - f_*}{\|w_0 - w_*\|^2}$ with L = 1 and $\mu = .01$. We compare

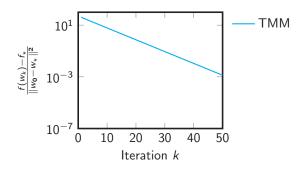
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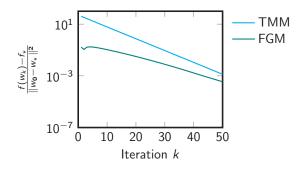
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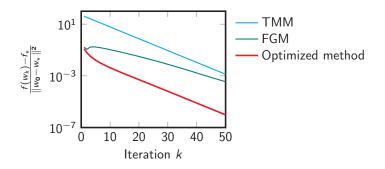
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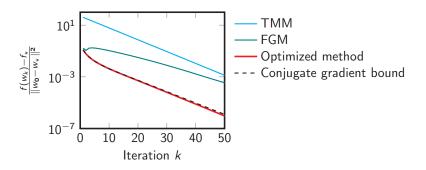
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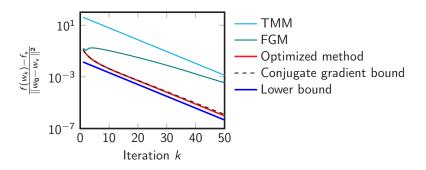
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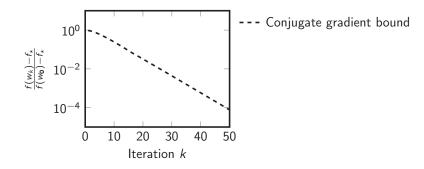
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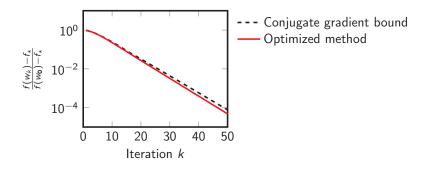
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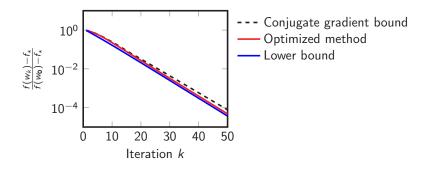
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Performance estimation problems

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Conclusions

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- ◊ overall: principled approach (definition of worst-case).

Difficulties:

 suffers from standard caveats of worst-case analyses, key is to find good assumptions/parametrization

Performance estimation's philosophy

- numerically allows obtaining tight bounds (rigorous baselines),
- $\diamond~$ results can only be improved by changing algorithm and/or assumptions,
- helps designing analytical proofs (reduces to linear combinations of inequalities), proofs can be engineered using numerics & symbolic computations!
- ◊ fast prototyping:

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- ◊ overall: principled approach (definition of worst-case).

Difficulties:

- suffers from standard caveats of worst-case analyses, key is to find good assumptions/parametrization
- $\diamond~$ closed-form solutions might be involved.

Take-home messages

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When it does: principled approach to worst-case analyses.

Thanks! Questions?

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AdrienTaylor/Performance-Estimation-Toolbox on Github